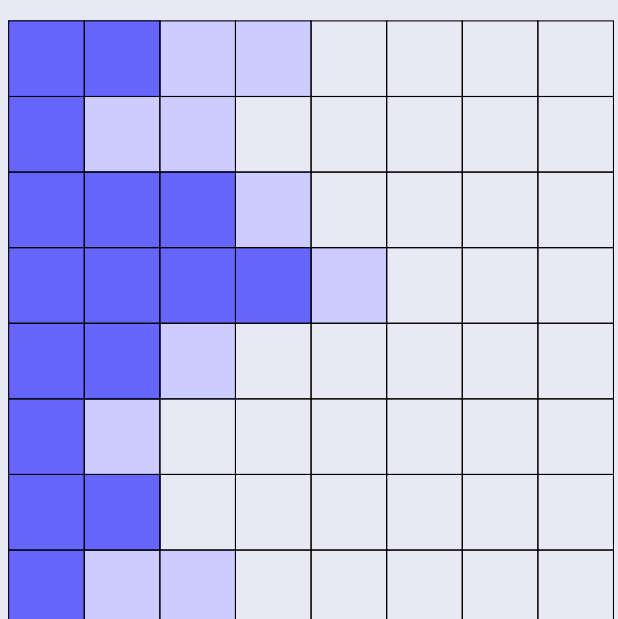
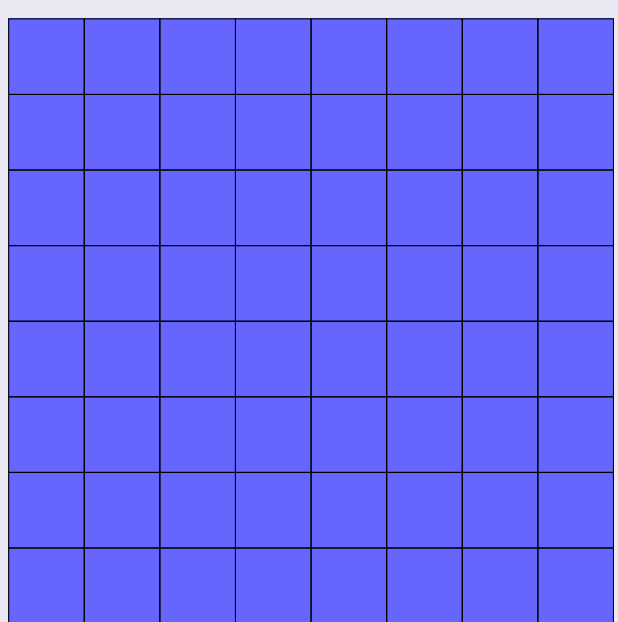


# Bounded Suboptimal Path Planning with Compressed Path Databases

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## Introduction



Compress First move matrix by:  
RLE, h-symbol

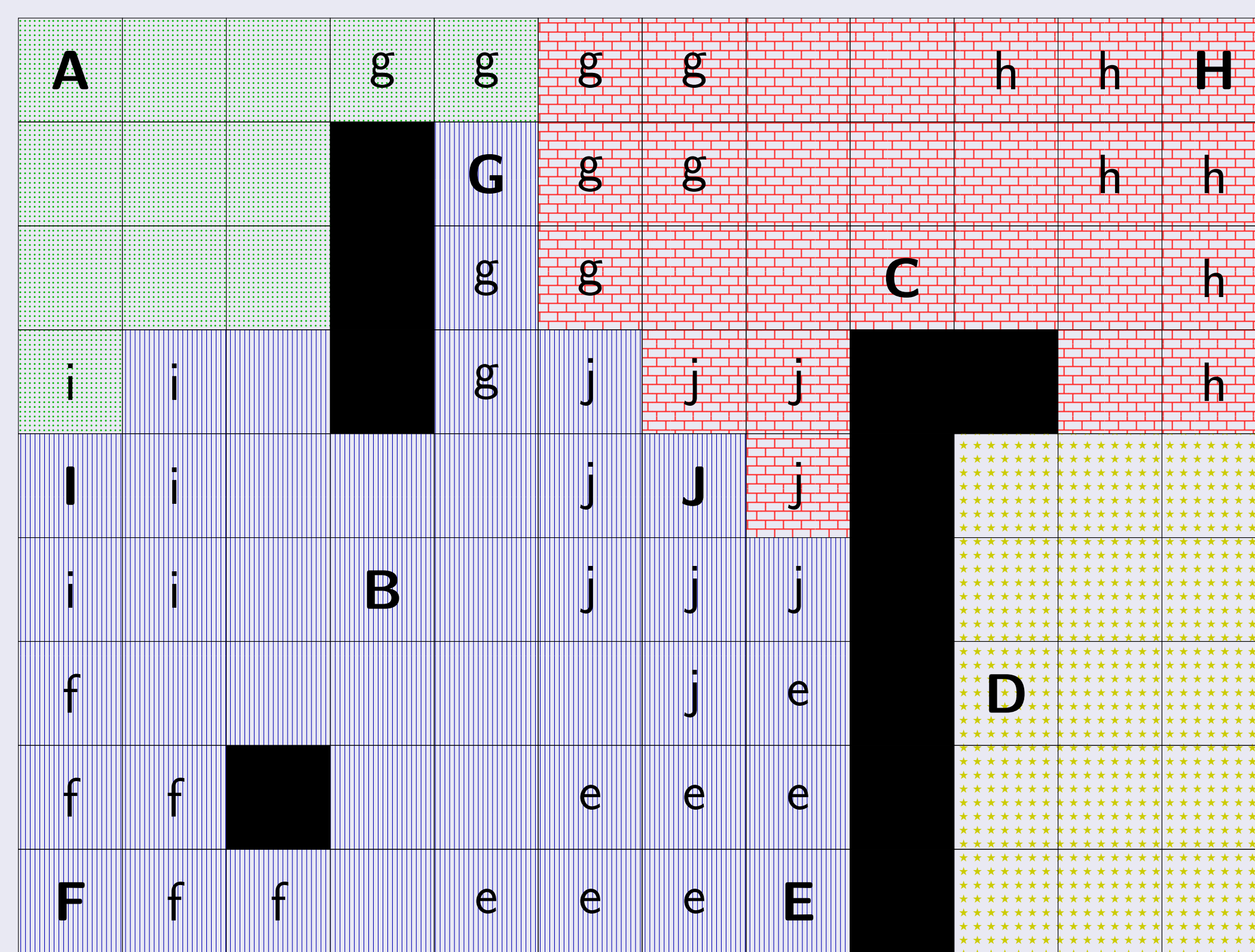
What are CPDs?

- Precompute all-pair first moves:  $O(n^2)$
- Compress:**  $\ll O(n^2)$
- Search-free path extraction

Motivation of this work:

- Existing approaches have achieved good compression and hard to improve.
- Trade-off between space and suboptimality.

## Method: Compute Centroids



- let  $\delta = 3$
- S1: generate seed-centroids by a heuristic function such that  $d(c_i, c_j) \leq 2\delta$
- S2: generate centroids on "borders"
- $\delta \leq d(c_i, c_j) \leq 2\delta$
- $|C_i| \geq \frac{\delta}{2}$
- thus  $i \leq \frac{2V}{\delta}$

## Idea

- For each node, only store first-move data for a subset of grid nodes  $C$ , called *centroid*
  - First-move matrix:  $N \times N \rightarrow N \times |C|$
- Each node  $i$  belongs to a centroid  $C(i)$
- Centroid path ( $cp$  for short)
 
$$cp(s, t) = shortestPath(s, C(t)) + shortestPath(C(t), t)$$
- Bounded suboptimal:  $|t, C(t)| \leq \delta$

## Method: Reverse CPD

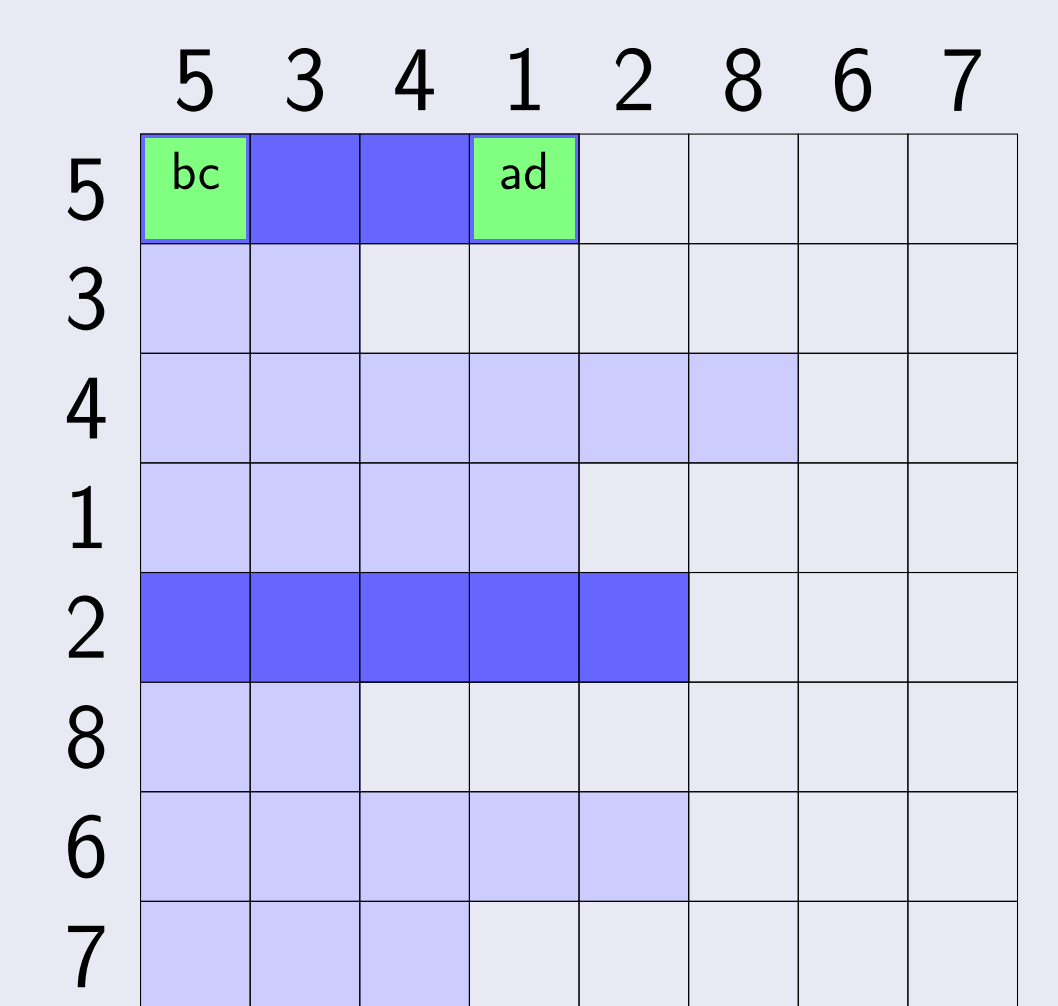
Advantages of reverse-CPD:

- get more space reduction from the centroids idea
  - forward-CPD throws non-centroids columns - may get same number of entries after compression
  - reverse-CPD throws non-centroids rows - always reduce the size
- faster path extraction
  - forward-CPD look-up in multiple row
  - reverse-CPD look-up in target row - less cache missing - single look-up may faster
  - and one look-up may extract multiple continuous moves - needs less number of look-up

Forward: 16, 10 (with centroids)



Reverse runs: 31, 9 (with centroids)



- $C = \{2, 5\}$
- Path: (2, 3, 4, 7, 5)

## Contribution

- Order-of-magnitude reductions in pre-processing
- Several factors improvement in database size
- Faster look-up
- Suboptimality cost per path is substantially smaller than the guaranteed upper bound

## Result

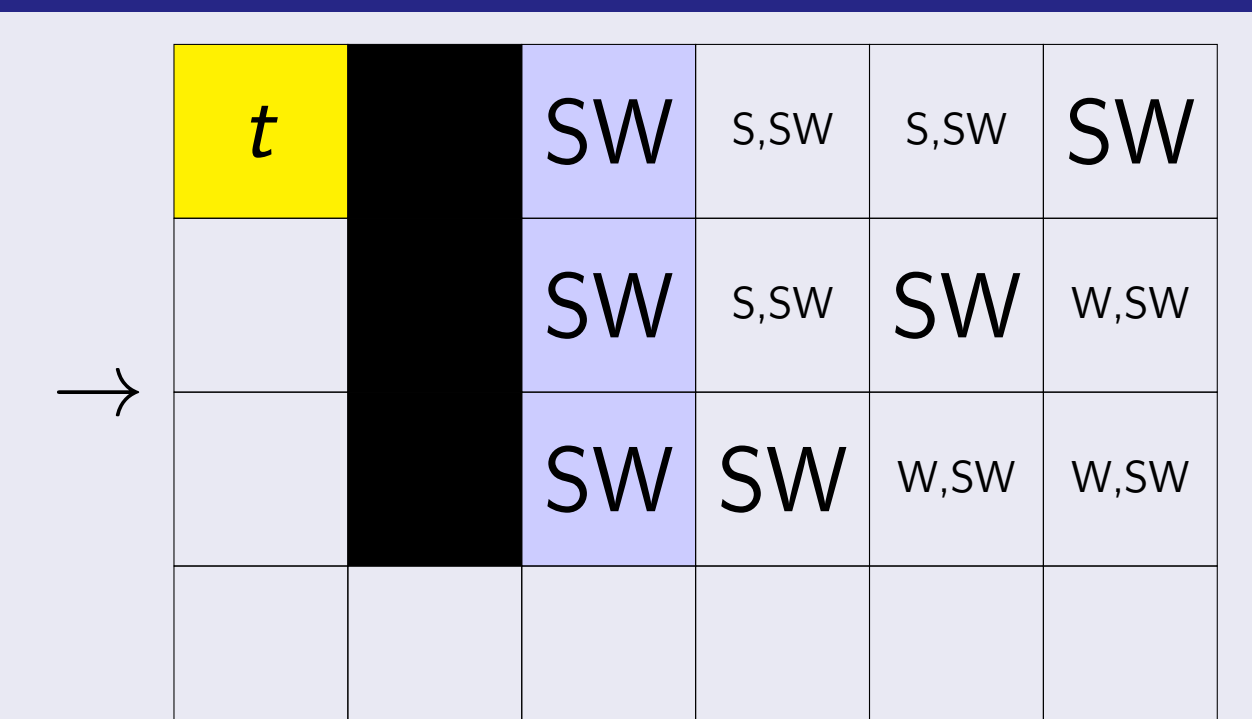
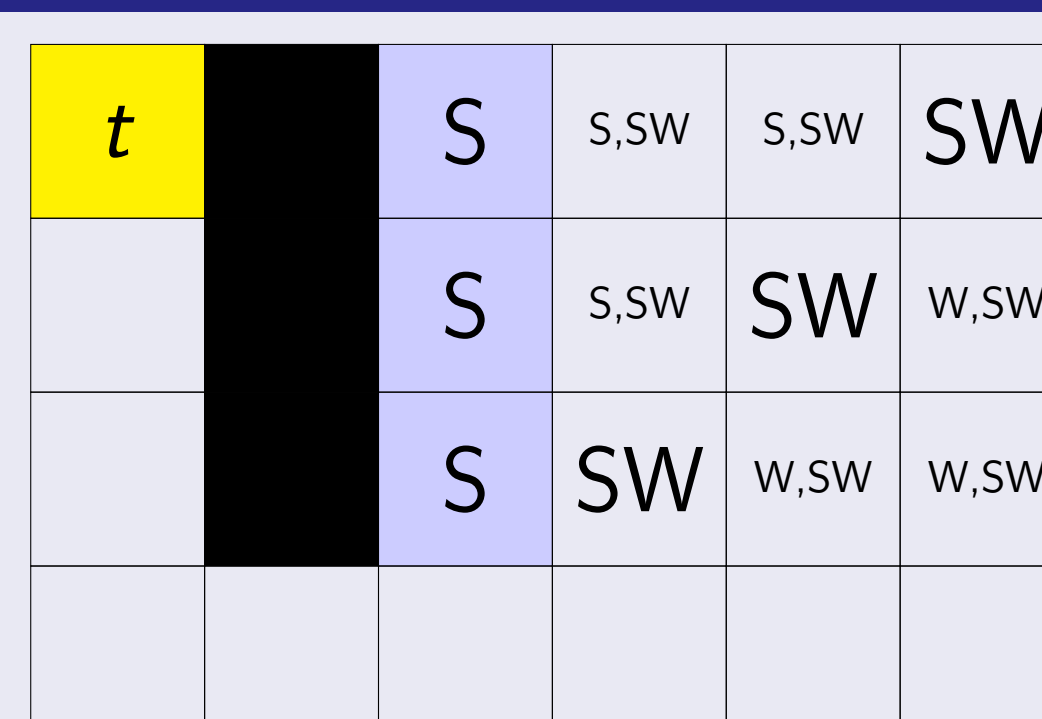
Size reduction ratio:  $\frac{|fwd_0|}{|rev_\delta|}$  or  $\frac{|fwd_0|}{|fwd_\delta|}$

$\delta$	type	mean	min	25%	50%	75%	max
2	rev	0.33	0.01	0.12	0.20	0.35	1.57
	fwd	0.97	0.85	0.93	0.98	1.01	1.19
4	rev	0.82	0.02	0.45	0.72	1.05	2.34
	fwd	1.13	0.85	1.02	1.12	1.24	1.68
8	rev	1.54	0.04	1.13	1.56	1.99	3.45
	fwd	1.34	0.86	1.08	1.27	1.56	2.90
16	rev	2.59	0.09	1.87	2.48	3.12	7.39
	fwd	1.60	0.86	1.14	1.36	1.86	4.90
32	rev	3.81	0.21	2.39	3.18	4.16	17.88
	fwd	1.84	0.87	1.19	1.42	2.10	7.20
64	rev	4.93	0.44	2.70	3.64	5.12	32.43
	fwd	2.08	0.87	1.24	1.50	2.21	9.65

Speed up ratio:  $\frac{Time(fwd_0)}{Time(rev_\delta)}$  or  $\frac{Time(fwd_0)}{Time(fwd_\delta)}$

	mean	min	25%	50%	75%	max
rev <sub>16</sub>	1.839	0.061	1.311	1.747	2.162	235.606
rev <sub>32</sub>	1.738	0.031	1.194	1.666	2.091	229.882
rev <sub>64</sub>	1.580	0.008	0.998	1.490	1.953	207.992
fwd <sub>16</sub>	1.209	0.013	1.038	1.125	1.312	175.824
fwd <sub>32</sub>	1.233	0.033	1.041	1.144	1.355	163.937
fwd <sub>64</sub>	1.230	0.012	1.013	1.139	1.389	184.361

## Method: Encode 'illegal' move



- Encoding  $S$  to  $SW$  allows us to compress the rectangle region to  $SW$ .
- When look-up column-3 symbols, we know  $SW$  is illegal, thus we can decode it to a valid move  $S$ .