



# Revisiting Bounded-Suboptimal Safe Interval Path Planning

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## 1. Problem Statement

**Environment** – a graph (grid with  $2^k$  connectedness or a roadmap)

**Agent** – an open disk of radii  $r$  that is allowed to:

- a) **move** with constant speed
- b) **stay** and **wait** any arbitrary amount of time
- c) **rotate** in place

**Dynamic obstacles** –  $N$  disks of radii  $r$  that move inside the environment along the known trajectories

**Trajectory** – a sequence of actions:  
 $\pi = \{a_1, a_2, \dots, a_k\}$

**Task** – having the given graph, a set of dynamic obstacles, start and goal agent's locations, find a trajectory that avoids static and dynamic obstacles.

Solution quality is determined by the cost(duration) of the agent's trajectory.

**Cost** – sum of actions' durations:  
 $c(\pi) = \sum_{i=1}^k dur(a_i)$

**Optimal Solution:**  
 $\pi^*$  = Plan with the minimal cost

**Bounded-suboptimal solution**  
 $\pi: c(\pi) \leq w \cdot c(\pi^*), w \geq 1.0$

## 2. Safe Interval Path Planning

**SIPP** [Phillips and Likhachev, 2011] – A\* in a different search space.

**SIPP is complete and optimal**

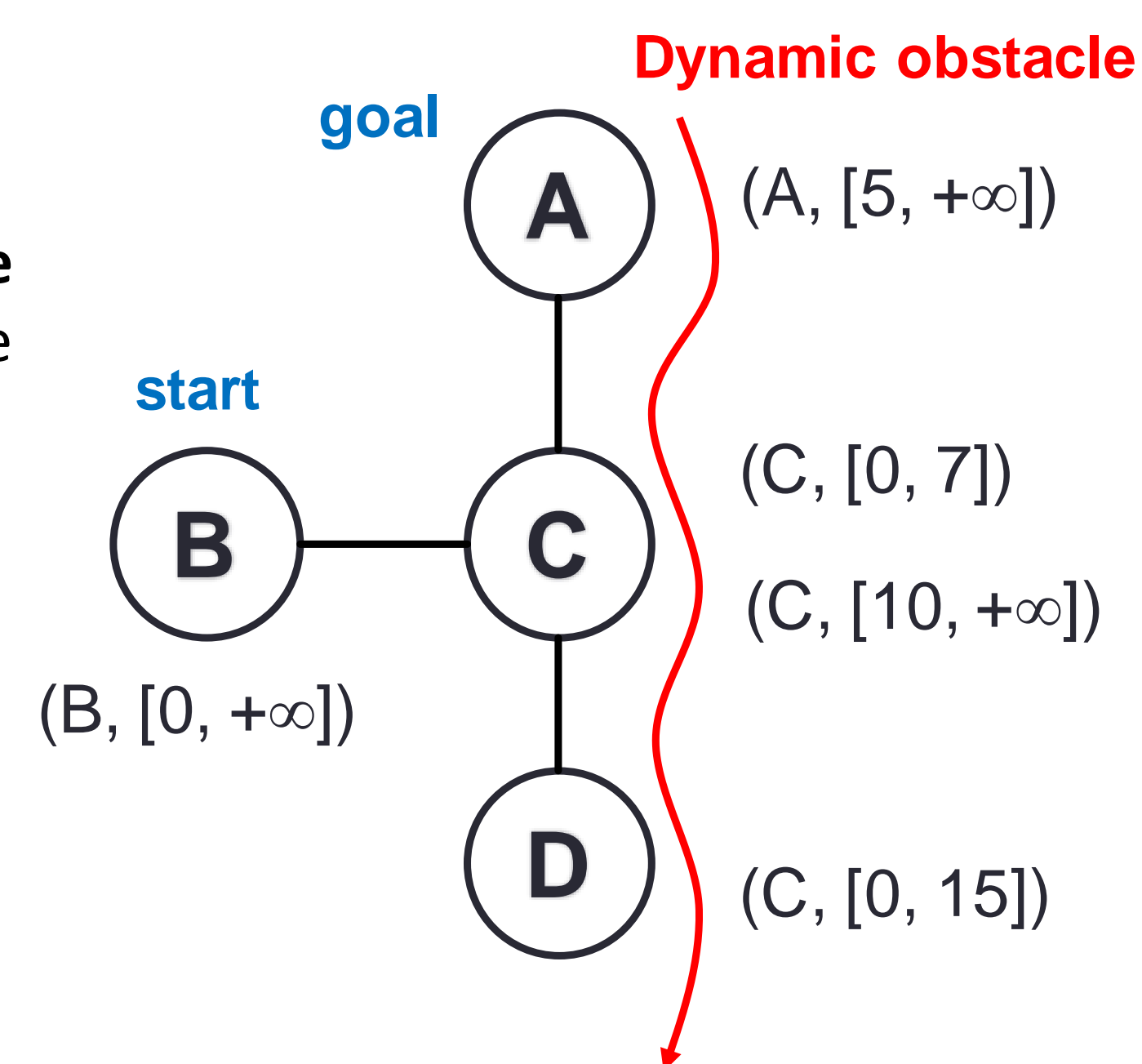
Each vertex has a corresponding **safe intervals** computed with respect to the trajectories of dynamic obstacles

**SIPP state** – (graph vertex, safe interval)  
 $n = (v, [t_i, t_j])$

$g(n)$  – the earliest possible arrival time to  $n$   
 $h(n)$  – estimate of the plan cost to the goal

- Admissible
- Consistent

When SIPP expands a node  $n = (v, [t_i, t_j])$ , then  $g(n)$  is the earliest possible arrival time to  $v$  in  $[t_i, t_j]$ .



## 3. Bounded Suboptimal SIPP

**Weighted SIPP with re-expansions (WSIPP<sub>r</sub>)**

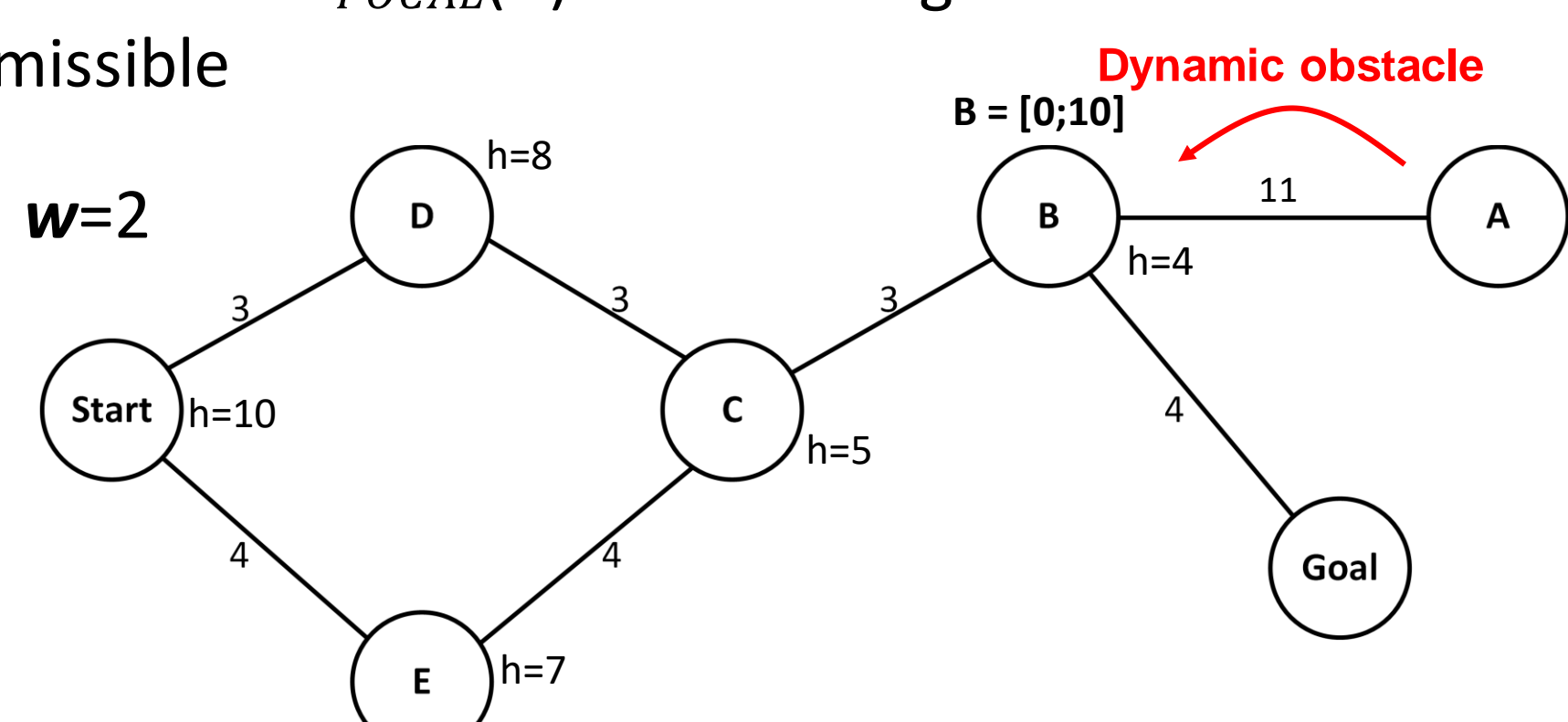
- h-values are multiplied by factor  $w$  (like in WA\*)
- Each state can be re-expanded multiple times
- Without allowed re-expansions loses the property of completeness (see example below)

**Weighted SIPP with Duplicate States (WSIPP<sub>d</sub>)**

- Suggested by authors of SIPP in [Narayanan et. al, 2012]
- Each successor is duplicated:  
Suboptimal version with f-value =  $g + w \cdot h$   
Optimal version with f-value =  $w \cdot (g + h)$
- Each successor is expanded at most once

**SIPP with FOCAL list (FocalSIPP)**

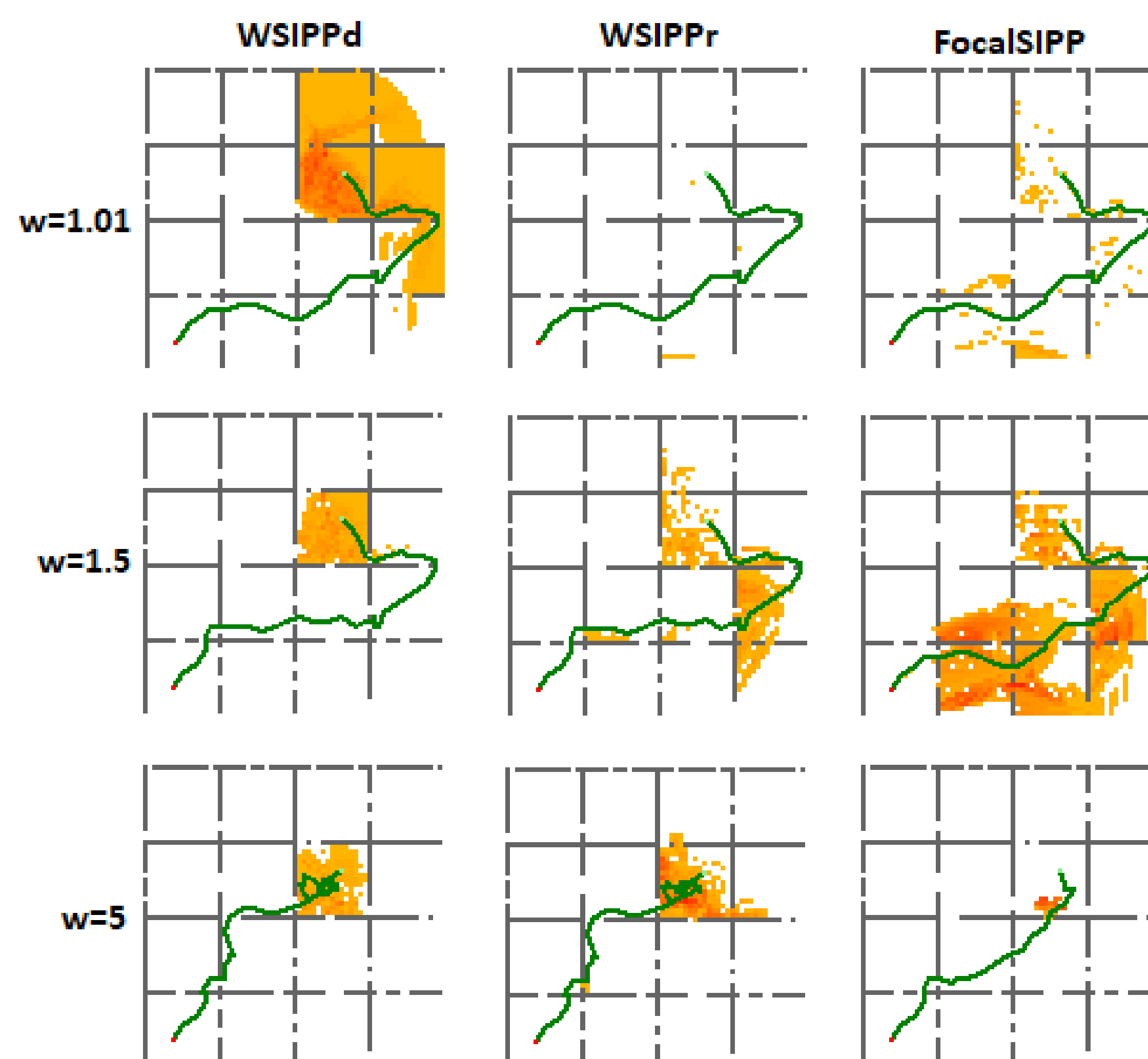
- Each state can be re-expanded multiple times
- Contains additional list – FOCAL [Pearl and Kim, 1982]
- FOCAL is a subset of OPEN list:  $\{n | n \in OPEN, g(n) + h(n) \leq w \cdot f_{min}\}$
- Auxiliary heuristic function  $h_{FOCAL}(n)$  for choosing a node from FOCAL does not have to be consistent, admissible



An example of instance where SIPP with weighted heuristic ( $w=2$ ) and disallowed re-expansions cannot find a solution. Safe-interval in B = [0;10]

## 4. Re-expansions Analysis

Re-expanded\* states are shown in yellow. The more states corresponding to the grid cell were re-expanded – the darker this cell is (remember that in SIPP numerous states corresponding to the same cell might exist). It is clearly seen that when  $w$  is close 1 the lowest number of re-expansions is achieved by **WSIPP<sub>r</sub>**, when  $w$  is large – by **FocalSIPP**, when  $w$  is "medium" – by **WSIPP<sub>d</sub>**



\*If **WSIPP<sub>d</sub>** expands both versions of some state, we consider it as a re-expansion

## 5. Experimental evaluation

**Maps:**

- Empty 64x64
- Warehouse 64x64
- Rooms 64x64
- Den520d 257x256

**Action model:**  
w/o (without) rotations  
w (with) rotations

**Grid connectivity:** 8, 16, 32

**250 dynamic obstacles**  
**100 instances per each map**

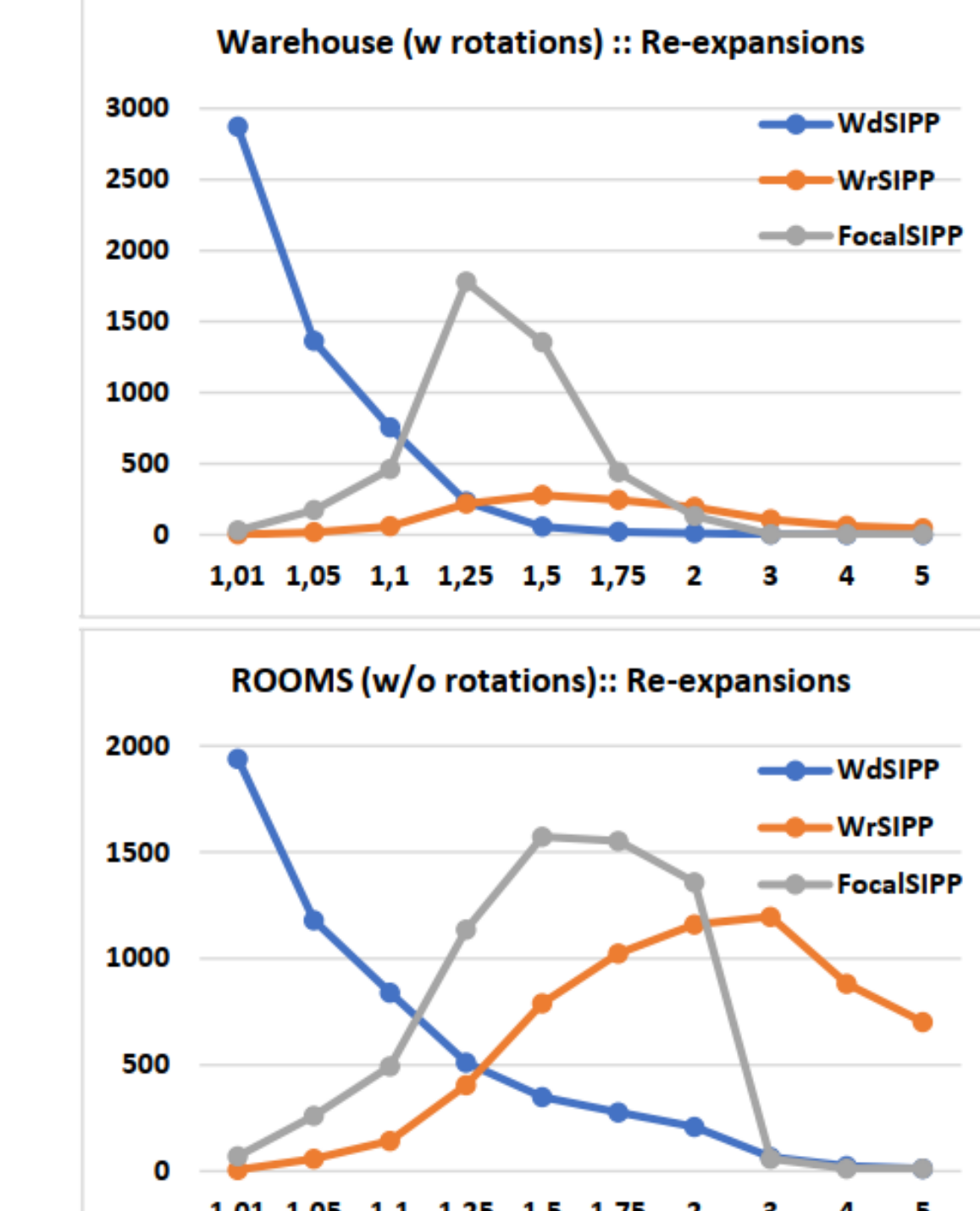
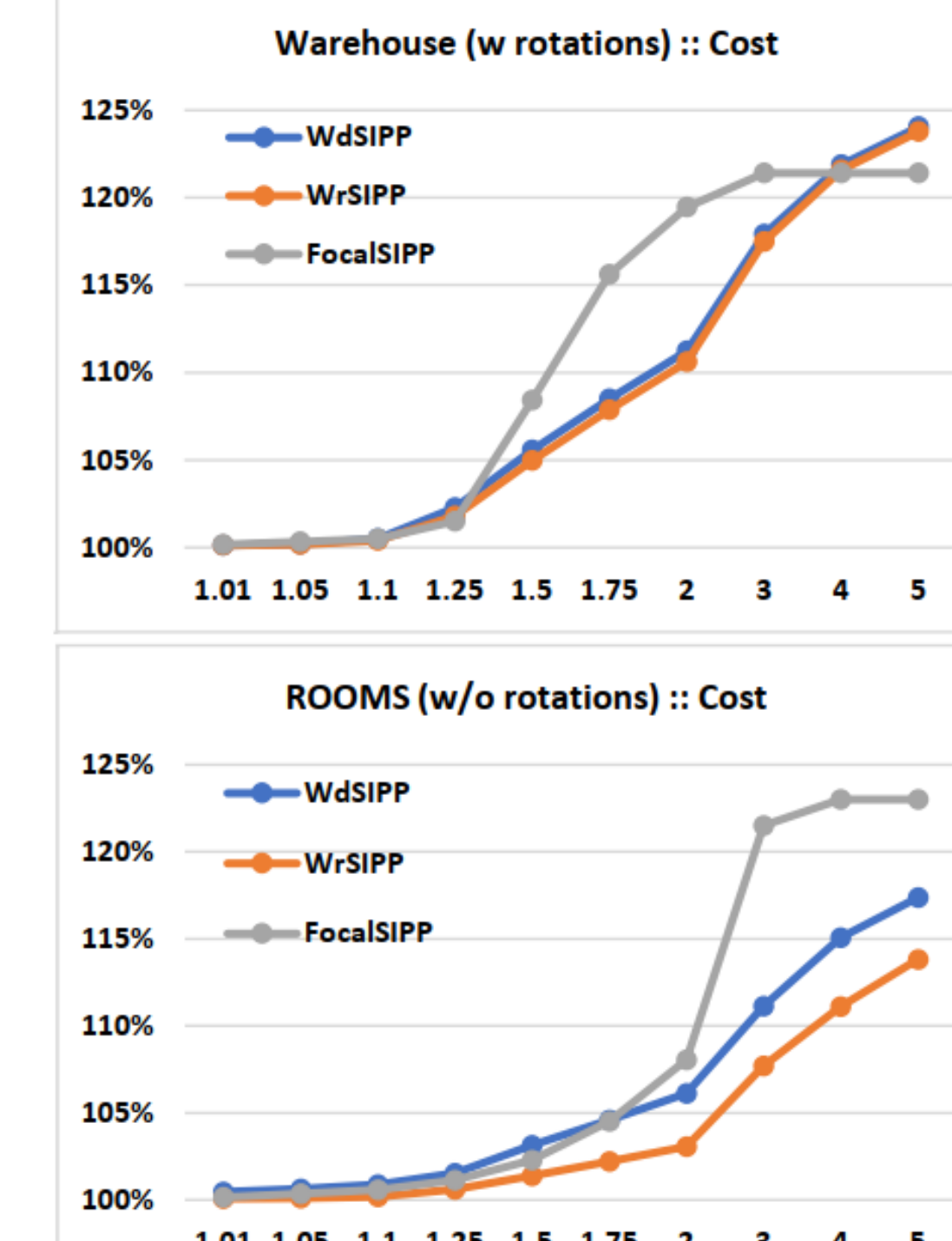
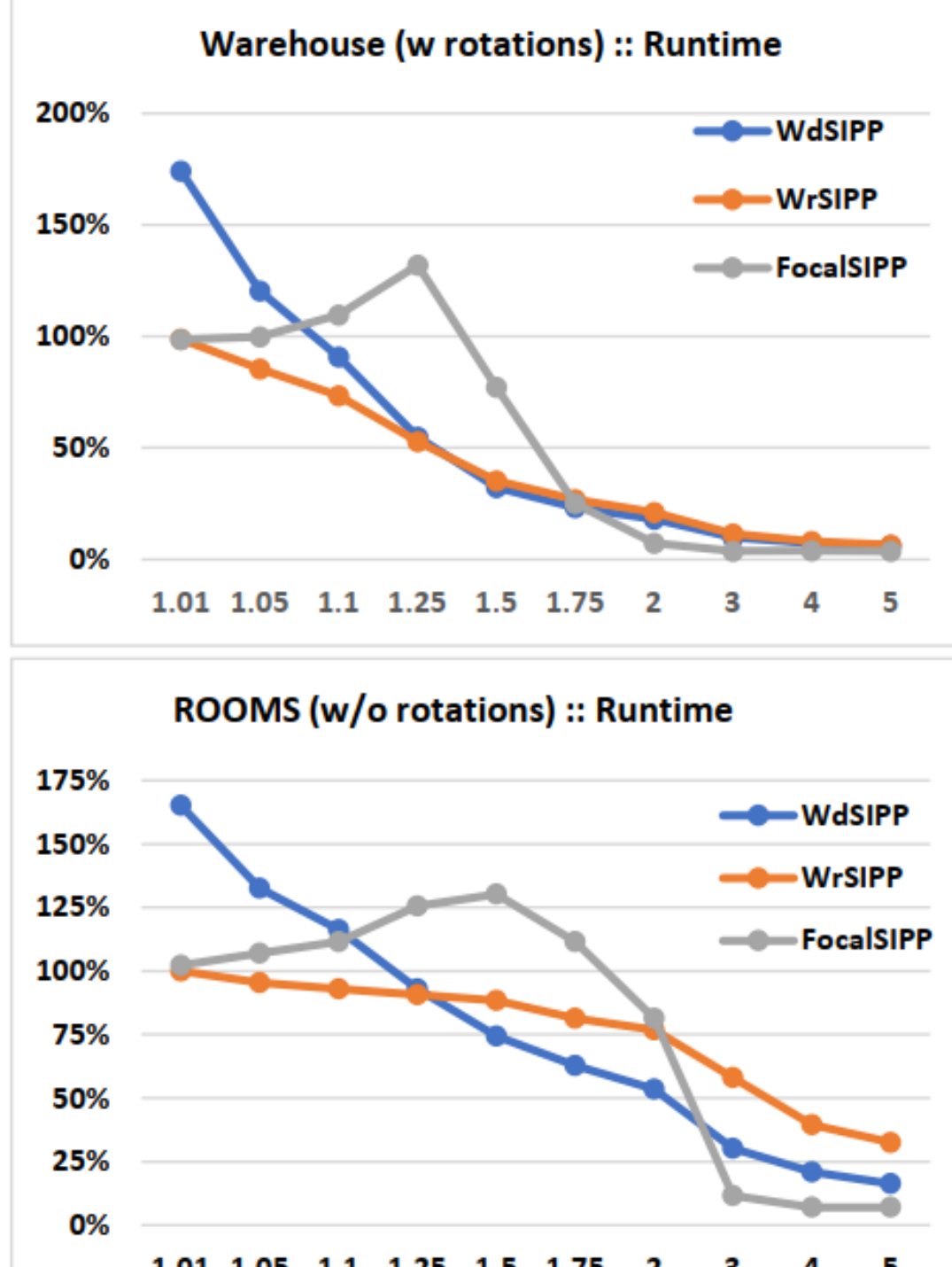
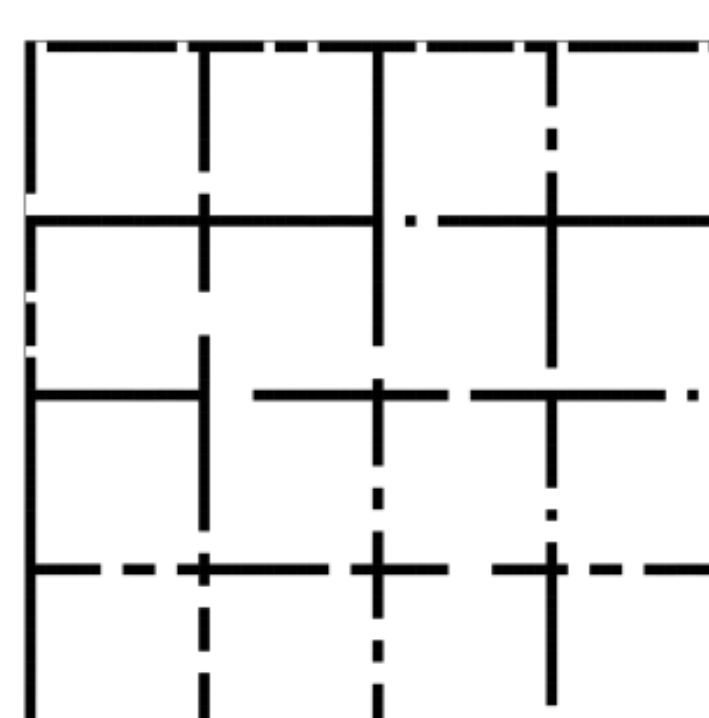
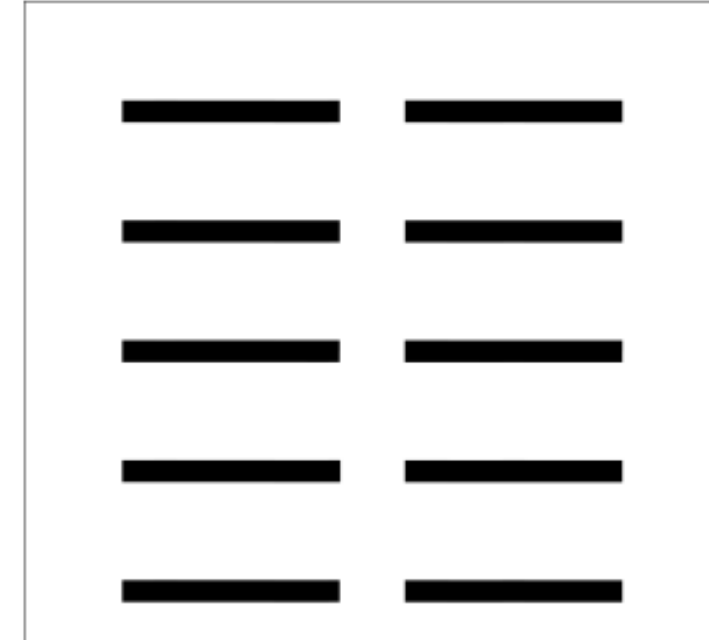
**Heuristic functions:**  
 $h$  – Euclidean distance;  
 $h_{FOCAL}$  – hops-to-goal

**Suboptimality factor:** 1.01, 1.05, 1.1, 1.25, 1.5, 1.75, 2, 3, 4, 5

8-connected

16-connected

32-connected



## References

- [Phillips and Likhachev, 2011] Phillips, M., and Likhachev, M. SIPP: Safe interval path planning for dynamic environments. In Proceedings of The 2011 IEEE International Conference on Robotics and Automation (ICRA 2011), 5628–5635.
- [Narayanan et. al, 2012] Narayanan, V. Phillips, M., and Likhachev, M. Anytime safe interval path planning for dynamic environments. In Proceedings of The 2012 IEEE/RSJ International Conference on Intelligent Robots and Systems (IROS 2012), 4708–4715.
- [Pearl and Kim, 1982] Pearl, J., and Kim, J. H. 1982. Studies in semi-admissible heuristics. IEEE transactions on pattern analysis and machine intelligence (4):392–399.



GitHub



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