

Analyzing and Avoiding Pathological Behavior in Parallel Best-First Search

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Parallel Best-First Search (BFS)

- BFS is used to solve graph search problems such as planning
- Many parallelization methods have been previously proposed
- Previous parallel A* methods have been experimentally shown to scale well

Method	BFS	Environment	Citation
HDA*	A*	distributed	Kishimoto et al. 2009
PBNF	A*, WA*	multi-core	Burns et al. 2010
KPBFS	WA*	multi-core	Vidal et al. 2010
HDGBFS, LE, LG	GBFS	distributed	Kuroiwa and Fukunaga 2019

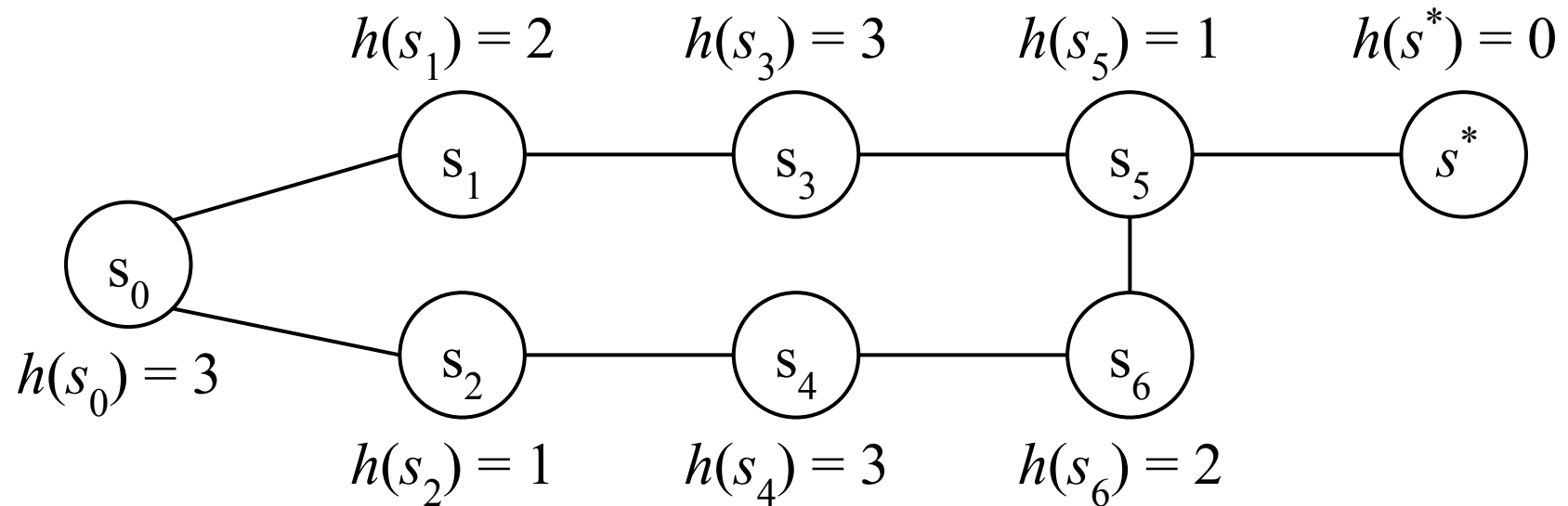
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- **What about parallel Greedy Best-First Search (GBFS)?**

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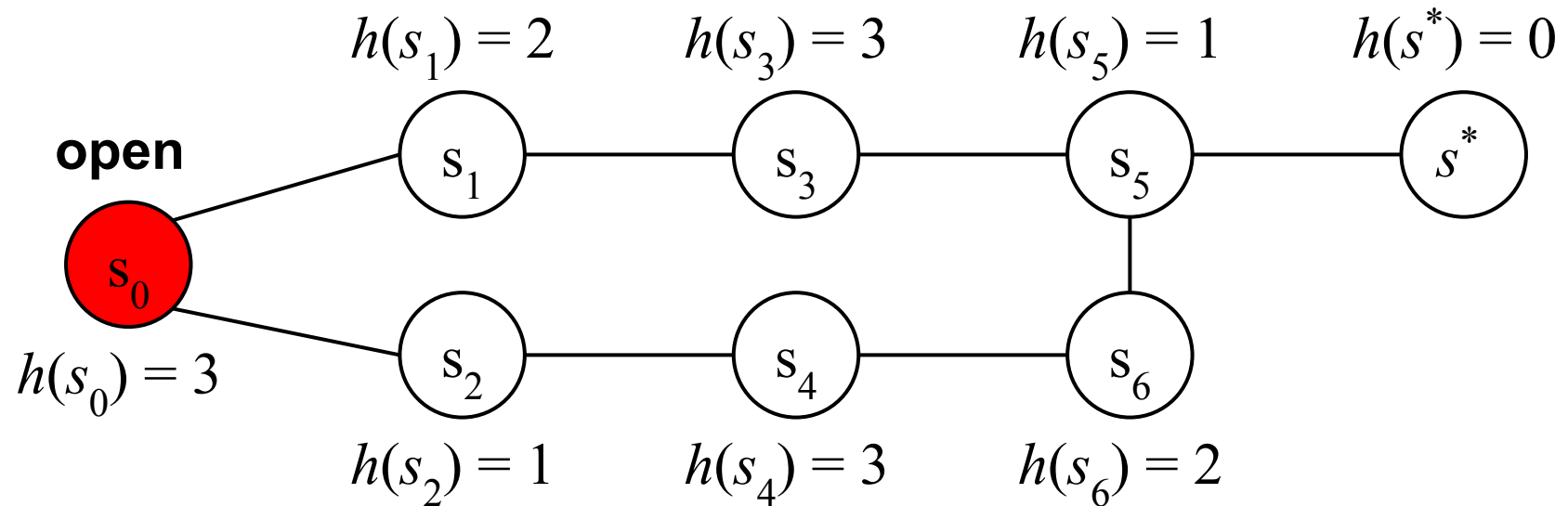
A Graph Search Problem Example with Greedy Best-First Search (GBFS)

- A solution of a graph search problem is a path from s_0 to s^*
- GBFS expands s with min. $h(s)$ (a heuristic value)



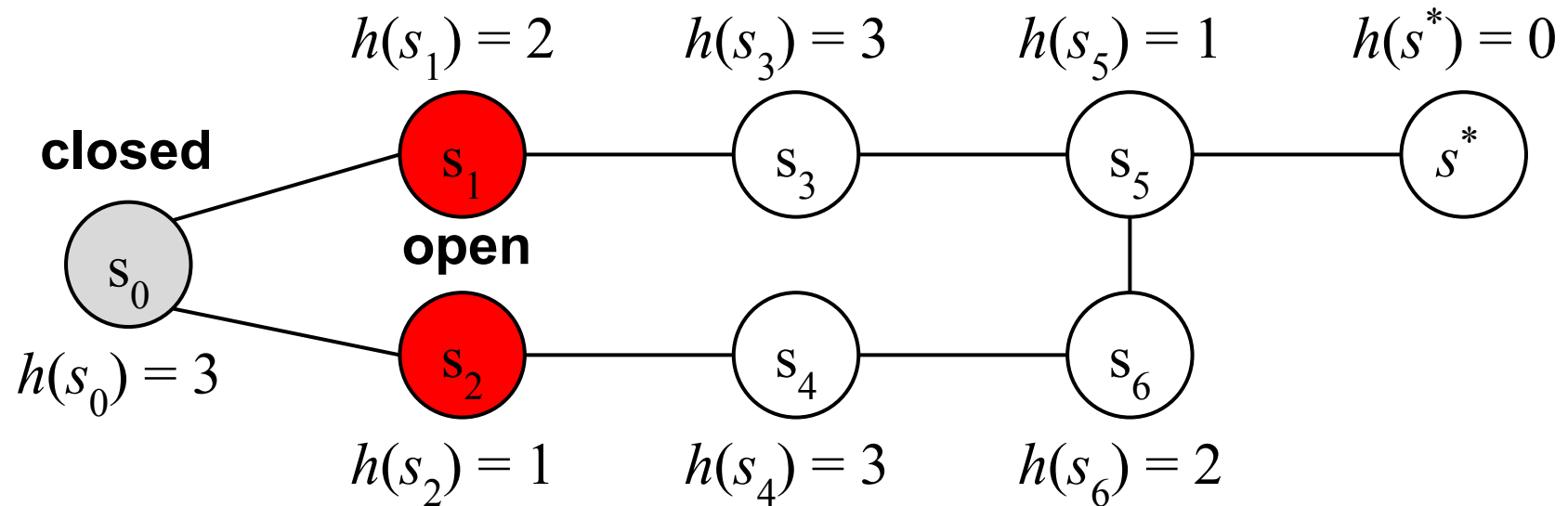
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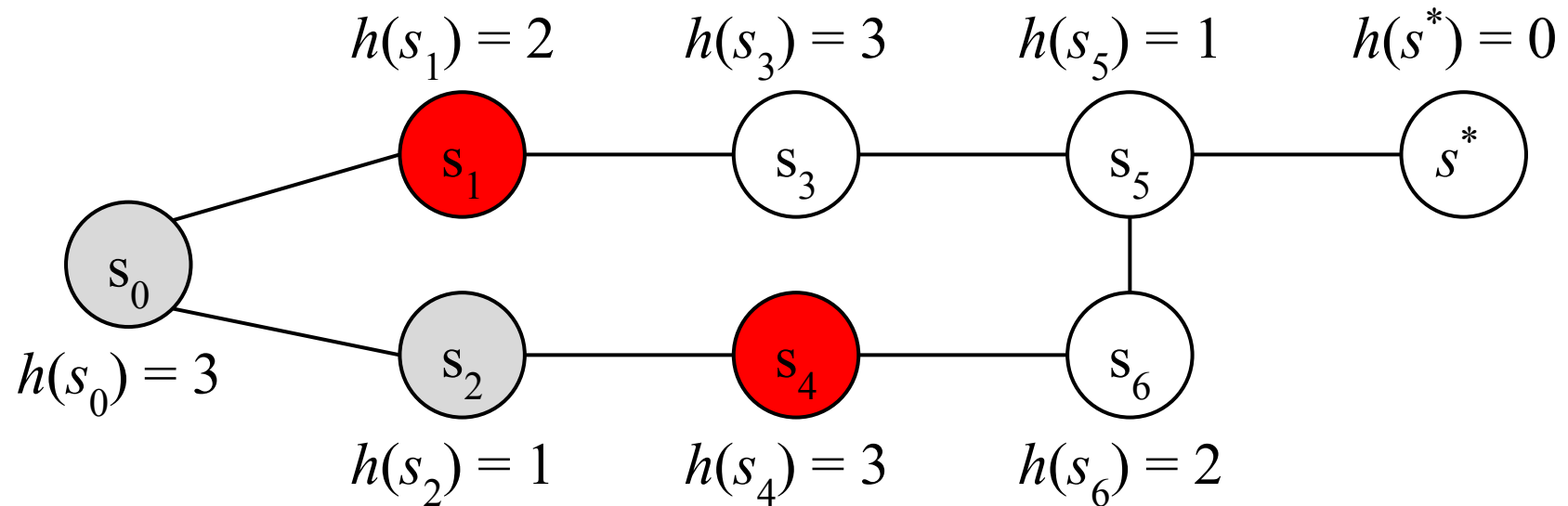
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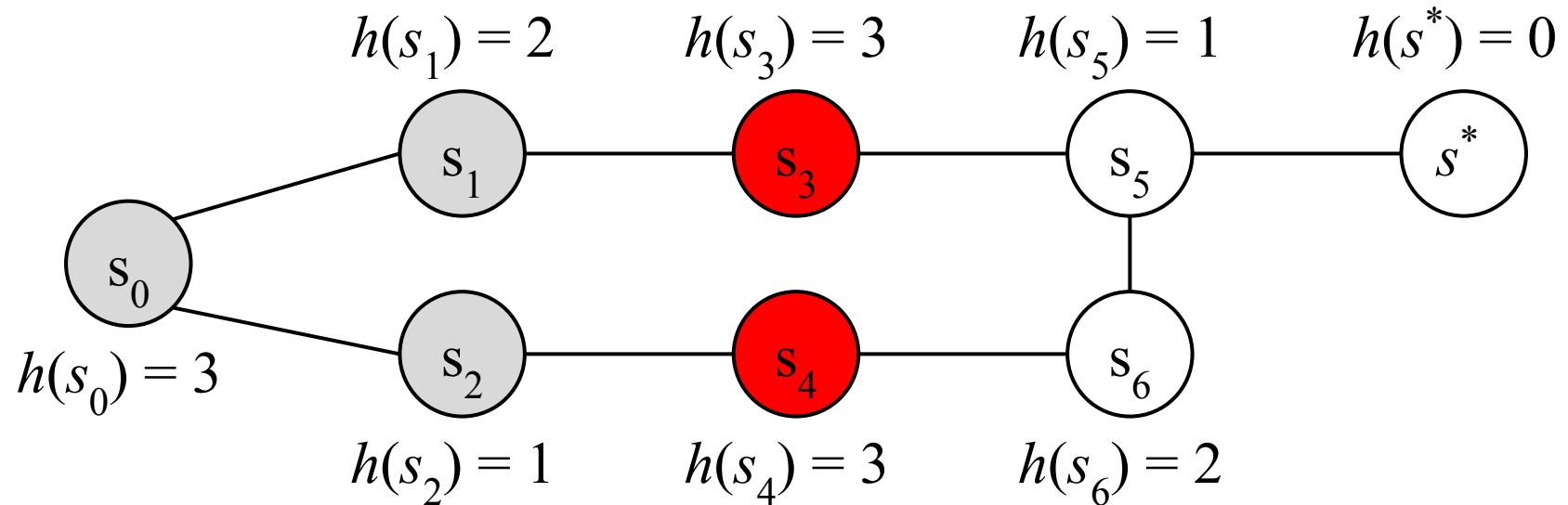
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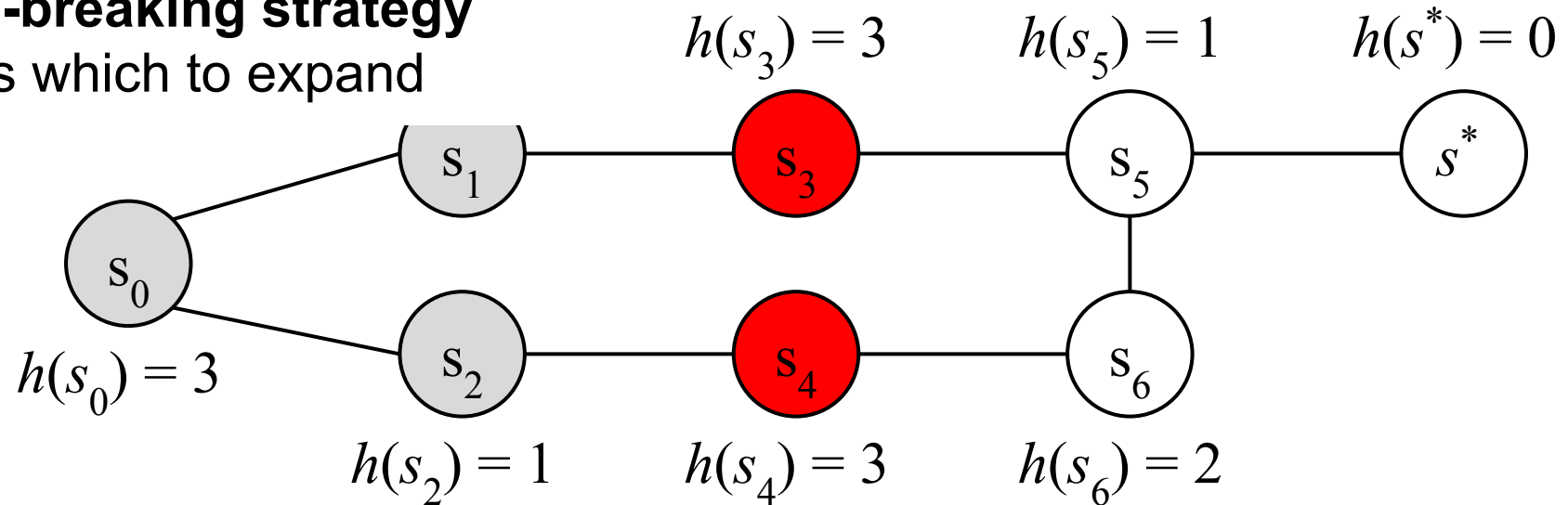
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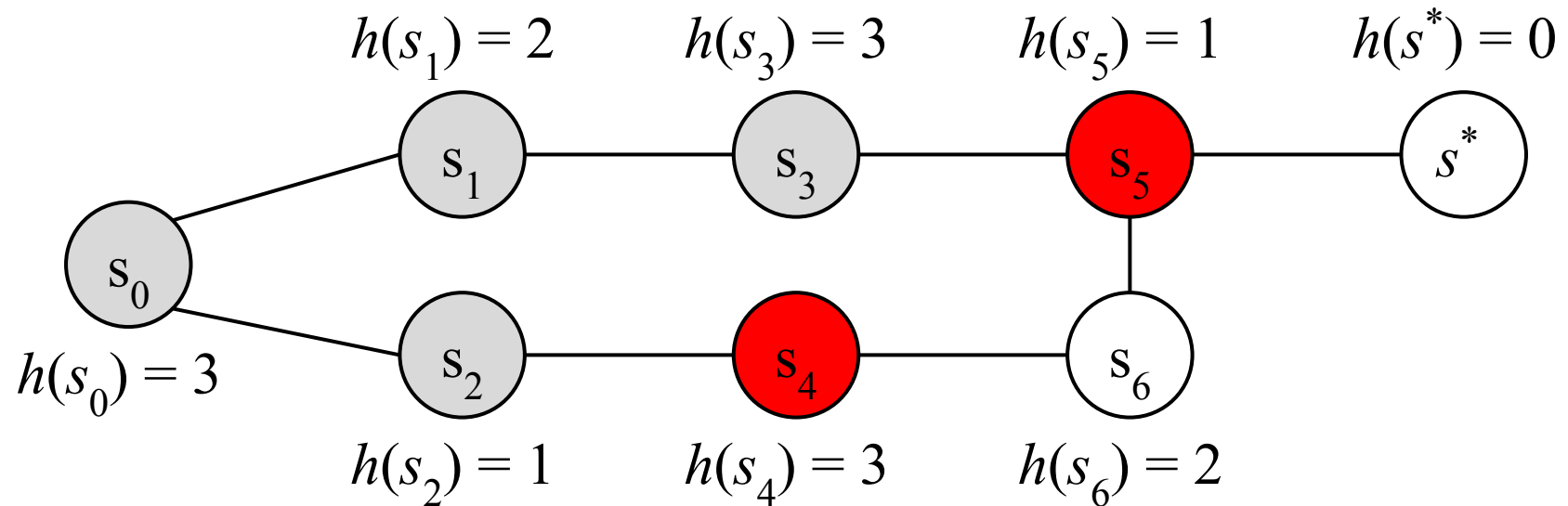
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The **tie-breaking strategy**
decides which to expand



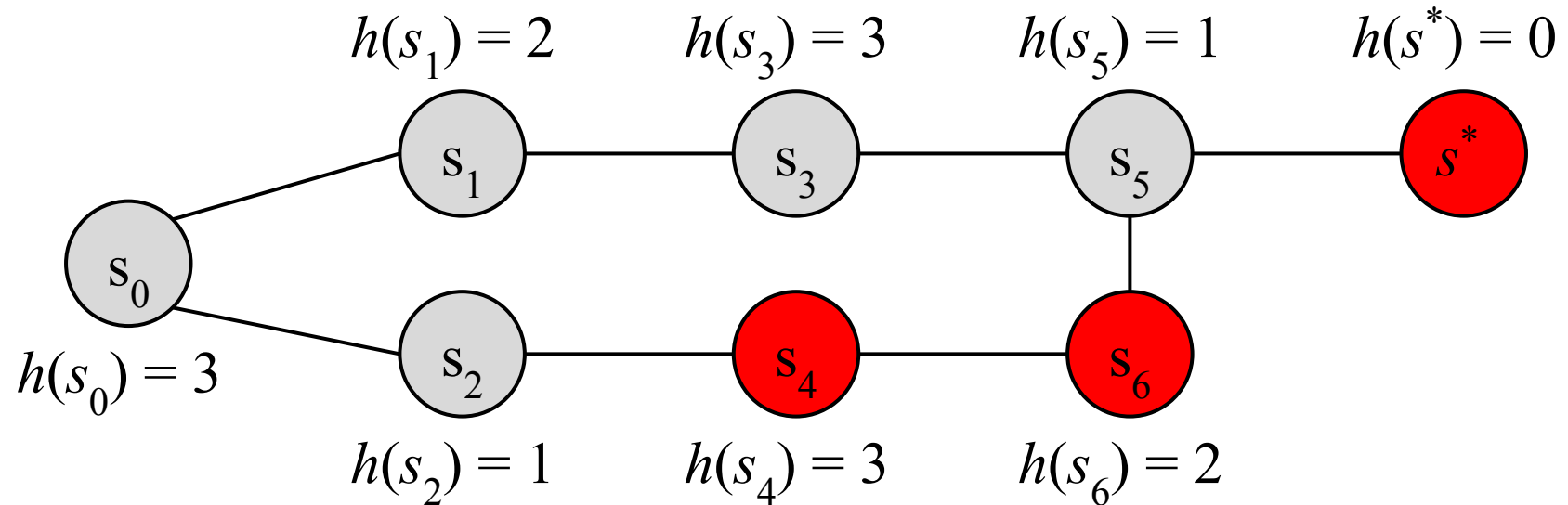
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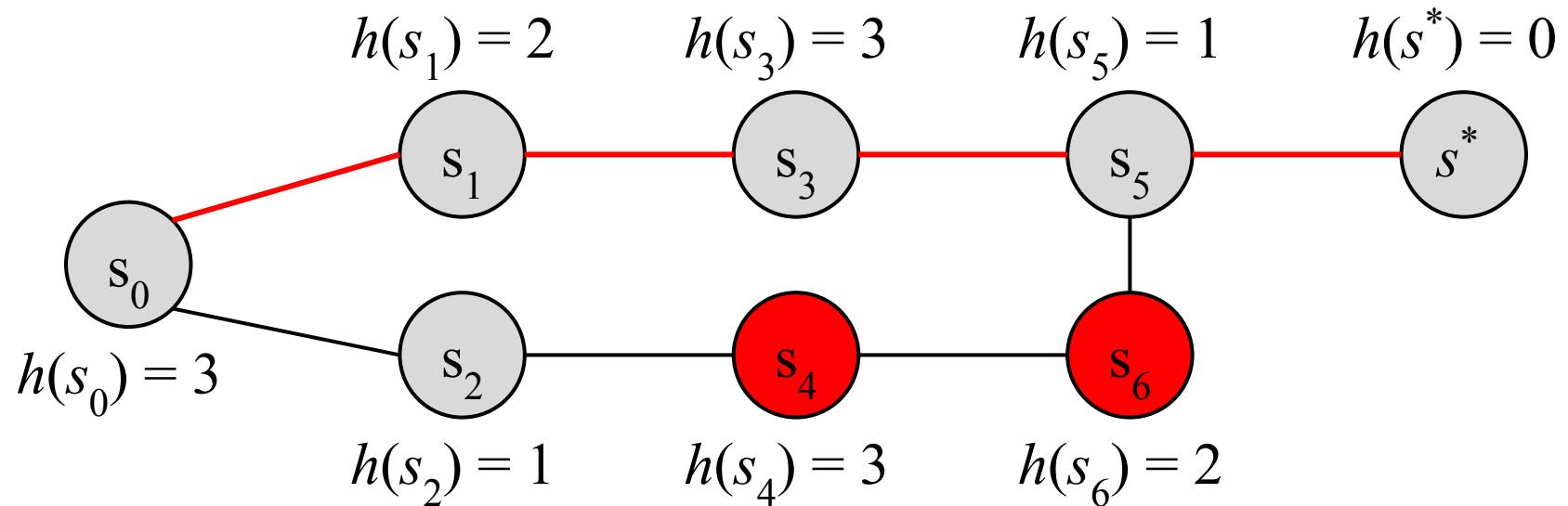
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Pathological Behavior in Parallel GBFS

[Kuroiwa and Fukunaga 2019]

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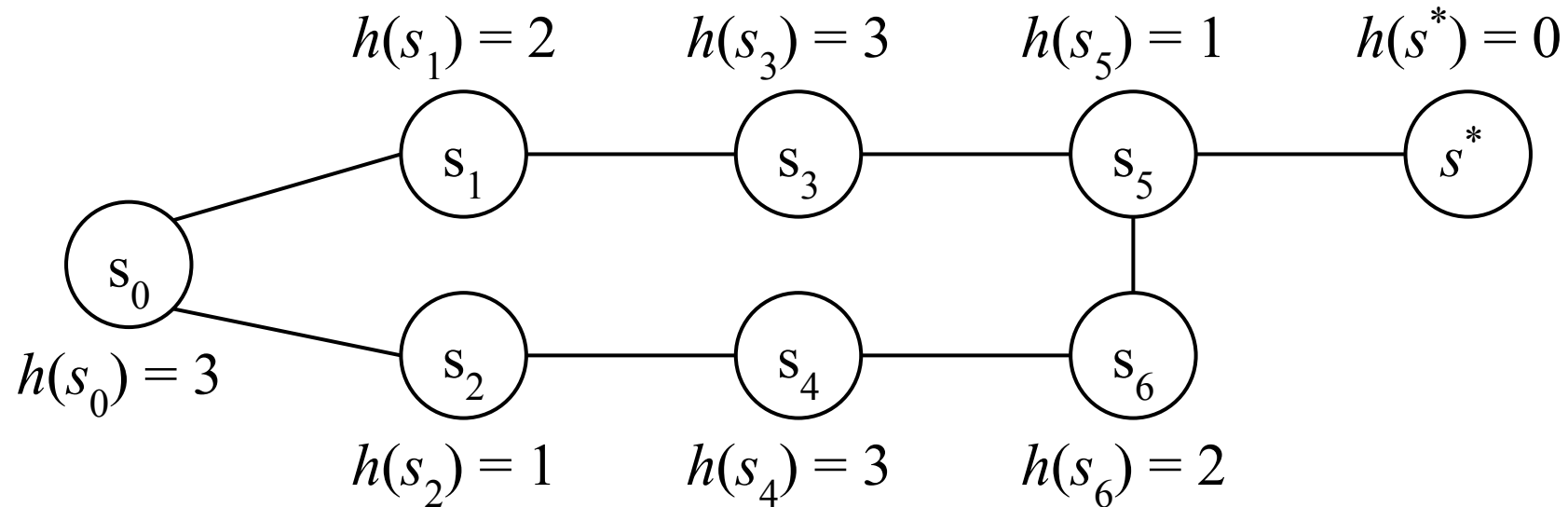
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- **Can we obtain theoretical bound for the performance degradation?**

KBFS: a Model of Parallel BFS

[Kuroiwa and Fukunaga 2019]

- KBFS [Felner et al. 2003]: similar to BFS, but simultaneously expands k states
- HDA*, KPBFBS, HDGBFS, and LE can be modeled as KBFS

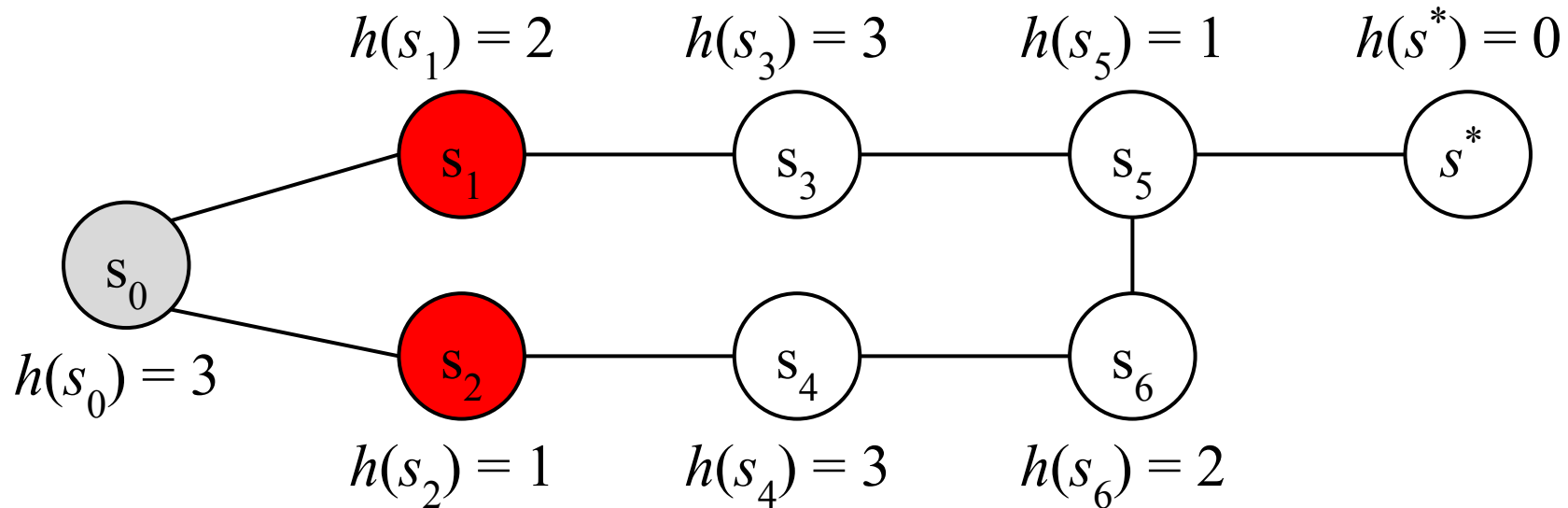


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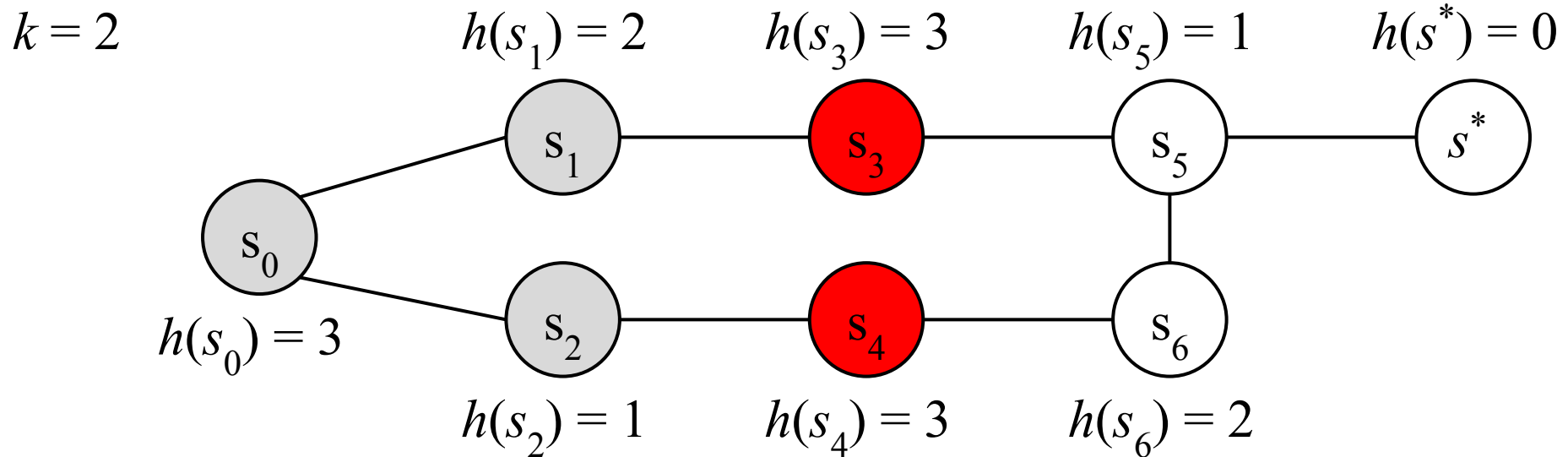
$k = 2$



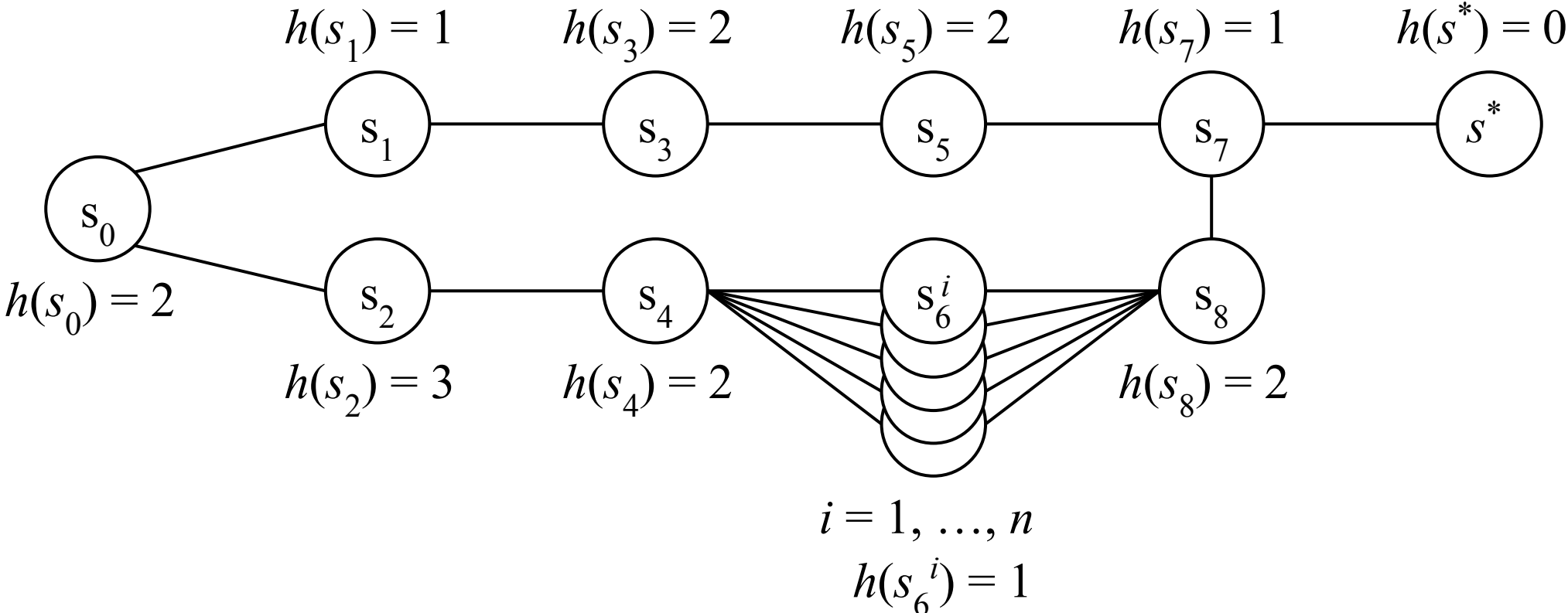
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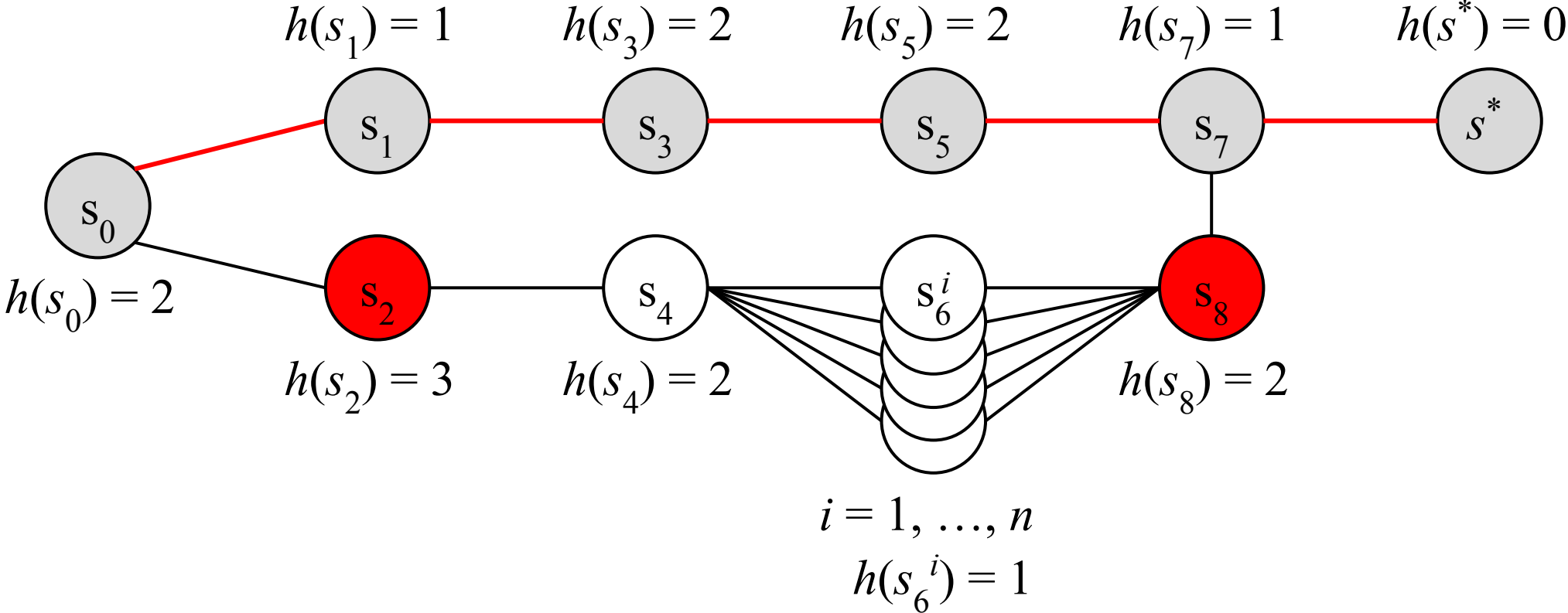


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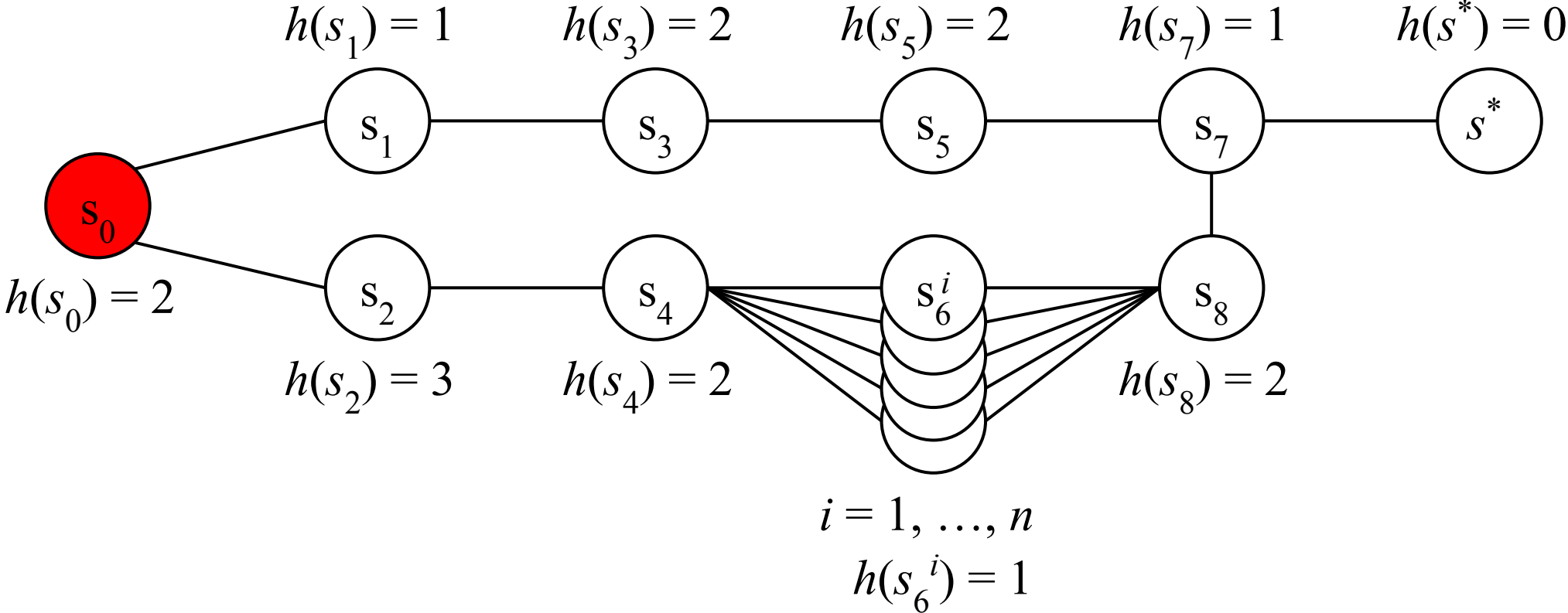
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GBFS expands 6 states



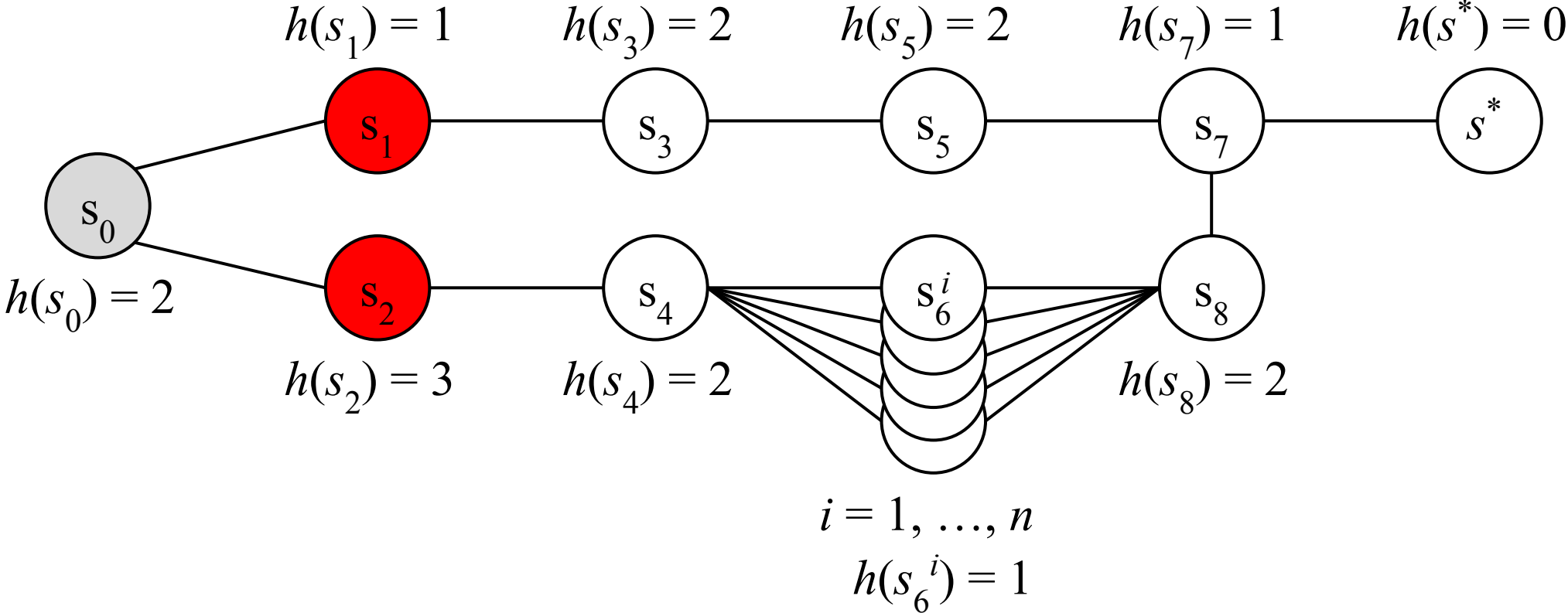
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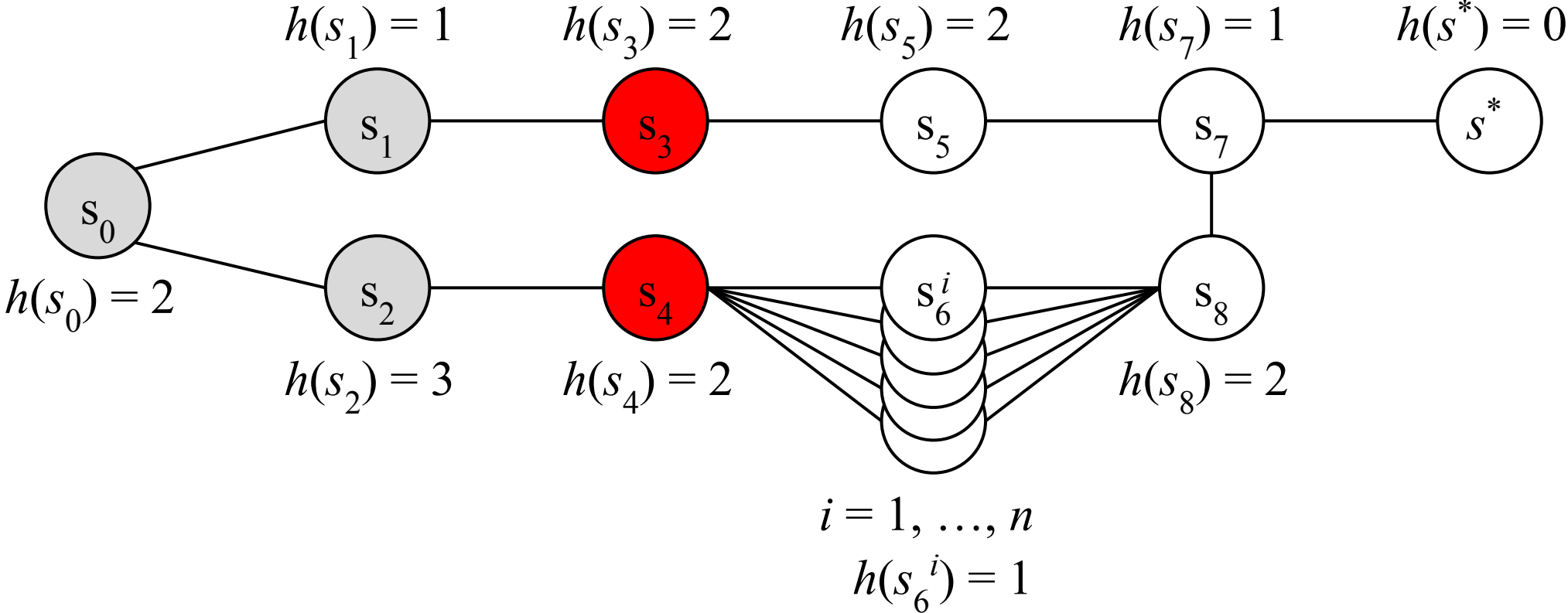
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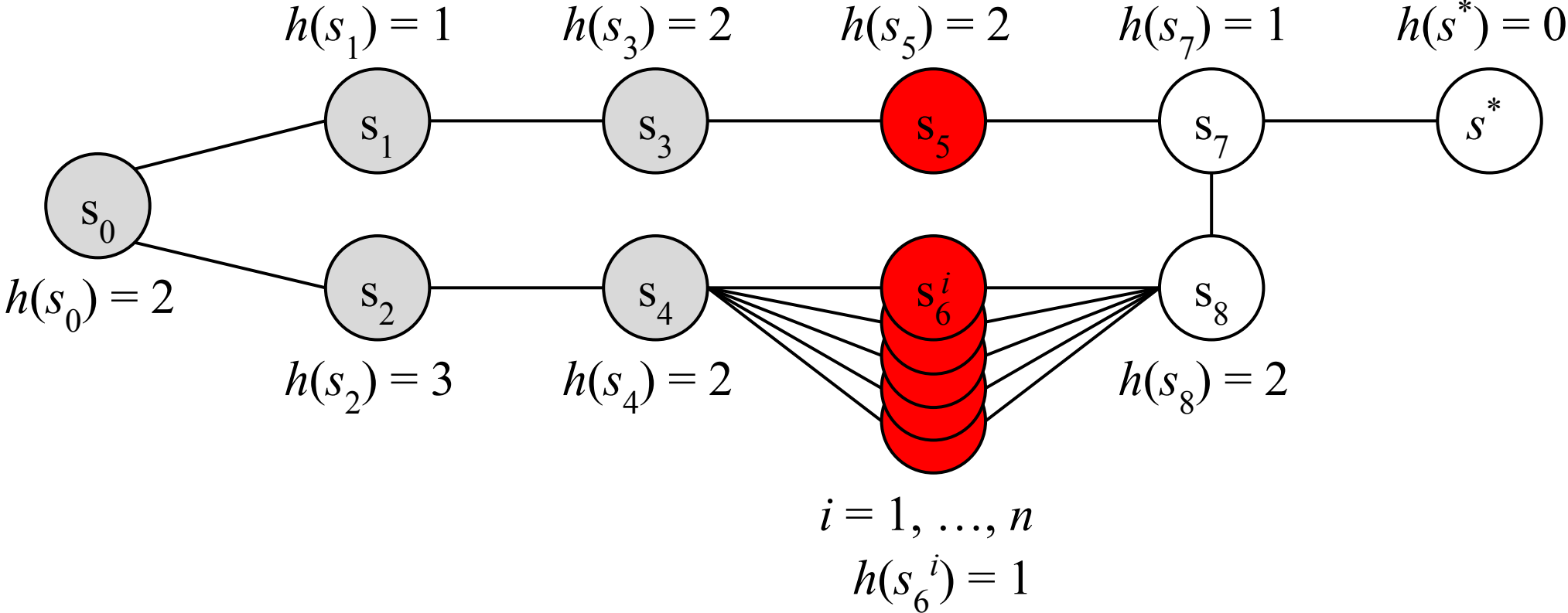
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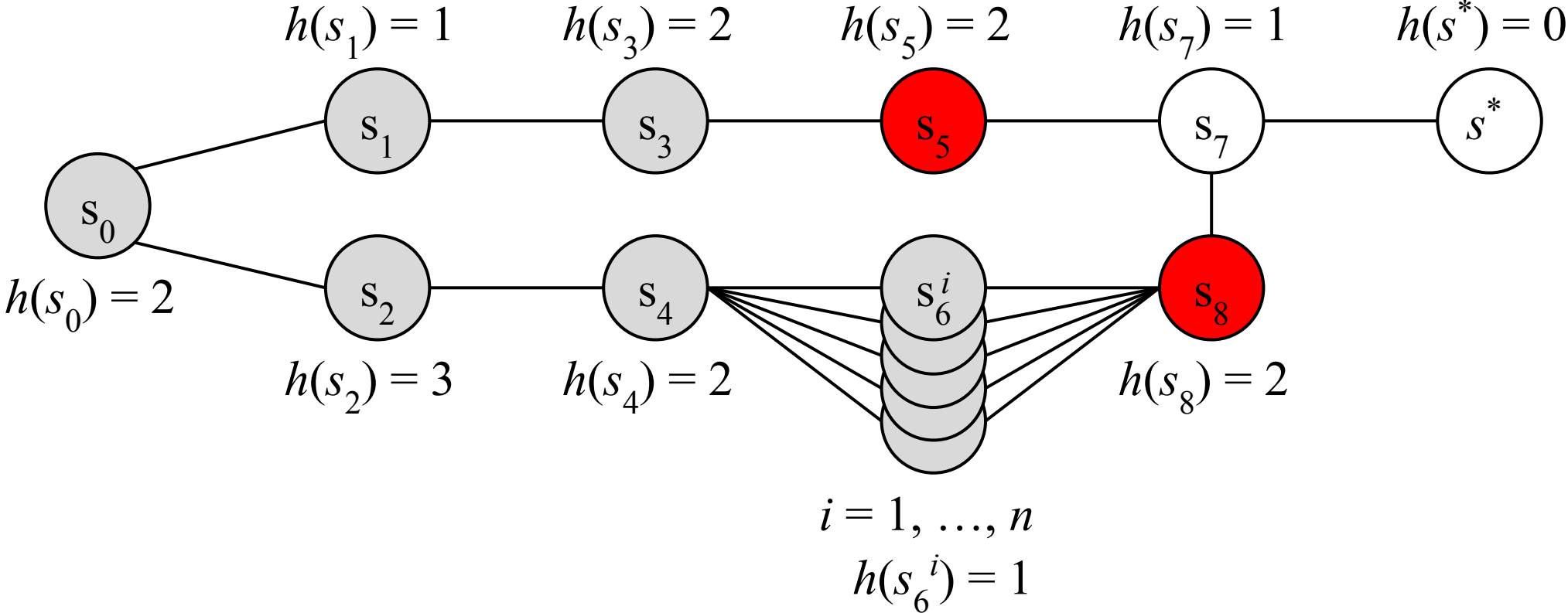
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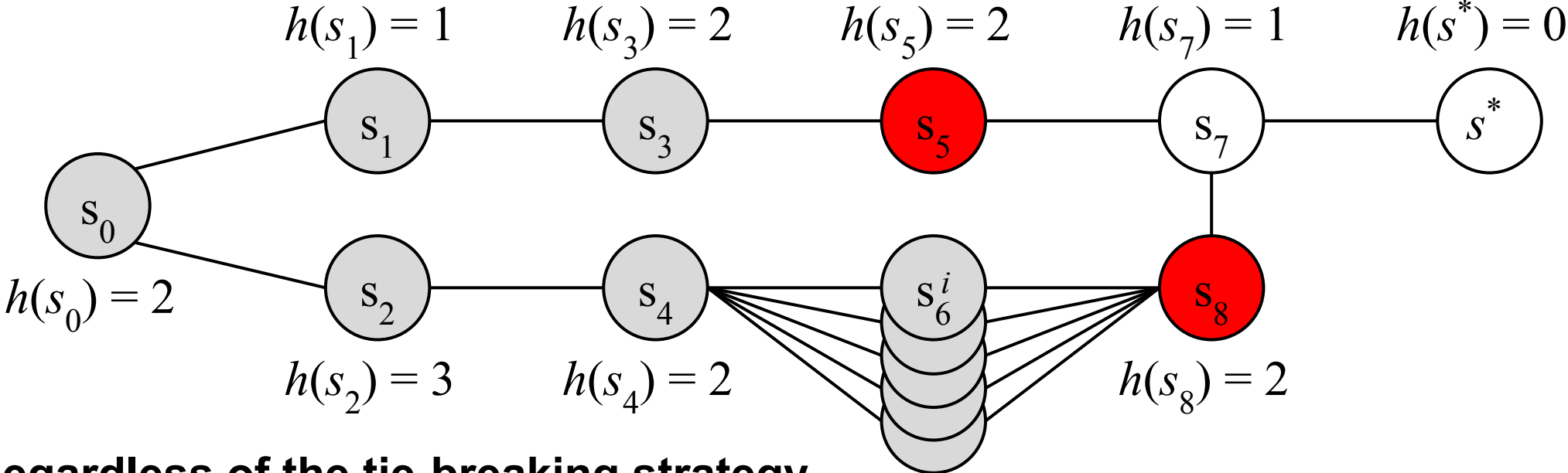
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Regardless of the tie-breaking strategy

$$i = 1, \dots, n$$

$$h(s_6^i) = 1$$

Pathology and t -Boundedness

- A is **pathological** relative to B if A expands arbitrarily more states than B does given some graph and heuristic
- A is **t -bounded** relative to B if A expands no more than t times as many states as B does, for any graph and heuristic

Pathologies in Parallel BFS

Method	Heuristic	t -boundedness	Model	BFS
HDA* (k processes)	consistent	k -bounded	KBFS (KA*)	A*
HDA*	inconsistent	pathological	KBFS (KA*)	A*
KPBFS ($w > 1$)		pathological	KBFS (KWA*)	WA*
KPBFS, HDGBFS, LE		pathological	KGBFS	GBFS
LG		pathological	KGBFS-like	GBFS

KA* is k -Bounded Relative to A* with a consistent heuristic

- A heuristic is **consistent**: $h(s') \geq h(s) + c(s, s')$
 $c(s, s')$ is the cost of the edge (s, s')
- KA* with any tie-breaking strategy is k -bounded relative to A* with **the worst-case tie-breaking strategy**

Proof Sketch: A* expands s with $\min. f(s) = g(s) + h(s)$ where $g(s)$ is the cost of the path from s_0 to s . A* expands each s with $f(s) < f^*$ (the cost of the optimal path). KA* expands s with $f(s) \leq f^*$ at every k expansions because s' on an optimal path has $f(s') \leq f^*$ by the consistency.

TB-Boundedness: Another Type of Bound

A is ***TB*-bounded** relative to *B* if *A* expands only states expanded by *B* with some tie-breaking strategy

Method	<i>TB</i> -boundedness	BFS
HDA*	no	A*
KPBFS ($w > 1$)	no	WA*
KPBFS, HDGBFS, LE, LG	no	GBFS

P_{GBFS}/C : a k -Bounded Parallel GBFS

- P_{GBFS} executes k threads of independent GBFS with different tie-breaking strategies in parallel [Kuroiwa and Fukunaga 2019]
Each GBFS does not change its expansion order
- P_{GBFS}/C uses the **shared evaluation cache** of heuristic values to speed up each GBFS while keeping the expansion order
- P_{GBFS} and P_{GBFS}/C are also *TB*-bounded

SPUHF: a *TB*-Bounded Parallel GBFS

- PUHF (Parallel Under High-water mark First) expands s only if $h(s) \leq h(\text{parent}(s))$ or any other thread is not expanding a state
- *Proof Sketch*: If $h(s) \leq h(\text{parent}(s))$, s is expanded by GBFS with some tie-breaking strategy
- SPUHF (Speculative PUHF) executes independent parallel search (**speculative search**) using idling threads with the **shared evaluation cache**

Bounded Parallel GBFS

- P_{GBFS}/C : k -bounded and TB -bounded
A parallel portfolio of independent GBFS with different tie-breaking strategies using the shared evaluation cache
 k -bounded and TB -bounded
- SPUHF: TB -bounded
A multi-core parallel GBFS similar to KPBFS
- See the paper for details

Experimental Results

Method	Coverage	# of solved instances unsolved by GBFS	# of unsolved instances solved by GBFS
LG	888	137	13
KPBFS	880	135	19
P_{GBFS}/C	928	164	0
SPUHF	864	115	15

Domain-wise results are complementary

55 domains from IPC-98-18
5 min. time limit
122 GiB memory limit

Conclusion/Summary of Contributions

- Proposed t -boundedness, Pathology, TB -boundedness
- Analyzed existing parallel BFS
- Proposed P_{GBFS}/C and SPUHF

Method	t -boundedness	TB-boundedness
HDA* [Kishimoto et al. 2009] (k processes), consistent heuristic	k -bounded	no
HDA*, inconsistent heuristic	pathological	no
KPBFS ($w > 1$) [Vidal et al. 2010]	pathological	no
KPBFS, HDGBFS, LE, LG [Kuroiwa and Fukunaga 2019]	pathological	no
P_{GBFS}/C [New]	k -bounded	TB -bounded
SPUHF [New]	unknown	TB -bounded